# Exact Ground States of Two-Dimensional $\pm J$ Ising Spin Glasses 

C. De Simone, ${ }^{1.2}$ M. Diehl, ${ }^{1}$ M. Jünger, ${ }^{1}$ P. Mutzel, ${ }^{3}$ G. Reinelt, ${ }^{4}$ and<br>G. Rinaldi ${ }^{5}$

Received January 20, 1996


#### Abstract

In this paper we study the problem of finding an exact ground state of a twodimensional $\pm J$ Ising spin glass on a square lattice with nearest neighbor interactions and periodic boundary conditions when there is a concentration $p$ of negative bonds, with $p$ ranging between 0.1 and 0.9 . With our exact algorithm we can determine ground states of grids of sizes up to $50 \times 50$ in a moderate amount of computation time (up to 1 hr each) for several values of $p$. For the ground-state energy of an infinite spin-glass system with $p=0.5$ we estimate $E_{0.5}^{\infty}=-1.4015 \pm 0.0008$. We report on extensive computational tests based on more than 22,000 experiments.


KEY WORDS: Branch and cut; Ising spin glasses; exact ground states.

## 1. INTRODUCTION

The last 20 years has witnessed a great deal of work on spin glasses. ${ }^{(6,13)}$ Nevertheless, the description of the phase transition and the nature of the ordered state remain controversial. ${ }^{(5,10)}$ The starting point for most theoretical work is the Edwards-Anderson (EA) model, whose Hamiltonian is

$$
\begin{equation*}
H(\omega)=-\sum_{i, j} J_{i j} \sigma_{i} \sigma_{j} \tag{1}
\end{equation*}
$$

[^0]where the $J_{i j}$ are random interactions and the spins lie on a regular lattice. An important case is the short-range model with Ising spins $\sigma_{i}$.

A widely studied model is the $\pm J$ Ising spin glass in which the sign of each bond is random but its magnitude is fixed. In two dimensions, the $\pm J$ spin glass with nearest neighbor interactions enters its spin glass phase only at zero temperature. ${ }^{(12)}$

In this paper we deal with two-dimensional models at zero temperature. We consider Ising spins $\sigma_{i}= \pm 1$ on an $L \times L$ lattice with periodic boundary conditions and nearest neighbor interactions. In this case, the total energy of the spin system is given by the Hamiltonian (1), with the sum restricted to pairs of nearest neighbor spins. Our goal is to compute an exact ground-state configuration of spin glasses when there is a concentration $p$ of negative bonds. As $p$ is increased from zero, one can observe a critical concentration $p_{c}$ which marks the phase transition between ferroand paramagnetism. Different publications give estimates for $p_{c}$ that lie between 0.09 and 0.163 . ${ }^{(2)}$

Since there are no partition functions in closed form that, given all interactions between the spins, yield a ground state, the only way to compute a ground state is using a numerical algorithm. However, the total number of states is $2^{L \times L}$, and so it is impossible, from a computational point of view, to find a ground state by just enumerating all possible states and computing the energy for each of them (unless $L$ is small enough). In fact, Barahona ${ }^{(1)}$ found a polynomial-time algorithm; however, his algorithm is also of high complexity and was never implemented, to the best of our knowledge.

Ground-state energies at $p=0.5$ have been calculated by several authors. Unfortunately, most of the methods proposed in the literature do not find the exact ground state, but an approximation for it. Such methods usually use Monte Carlo simulation ${ }^{(18)}$ and evolutionary ${ }^{(10.11)}$ and genetic algorithms; ${ }^{(15.16)}$ in ref. 4 a gauge-invariant method is proposed. In ref. 17 an exact integer method is proposed to find ground states for small systems.

The branch-and-cut method we proposed and described in ref. 3, which computes an exact ground state by finding a cut of maximum weight in a weighted graph, is able to solve large samples. It is the same method we used to find ground states for two-dimensional Ising spin glasses with periodic boundary conditions and nearest neighbor interactions based on a Gaussian bond distribution and an exterior magnetic field. ${ }^{(8)}$ With our method we can compute exact ground states of $\pm J$ Ising spin glasses on square lattices of size up to $L=50$ within 1 hr and up to $L=70$ within 1 day of computation time. We used $p \in\{0.1,0.3,0.5,0.7,0.9\}$. Because of the moderate computation time, we can give results based on more than 22,000 samples.

In the following section, we discuss the results we obtained for twodimensional systems.

## 2. COMPUTATIONAL EXPERIMENTS

Our computational experiments were carried out on a SUN SPARCstation 10. The complete computer code, except the linear programming routines, was written by us; to solve the linear programs we used the CPLEX Callable Library. ${ }^{(7)}$

Our experiments were done on $\pm J$ Ising spin glasses on $L \times L$ square lattices with a concentration $p$ of negative bonds, for a range of values of $L$ up to 50 . Indeed we can solve lattices of bigger size, say $70 \times 70$, but the running times for these systems are so high that we cannot run a sufficiently large number of these systems to get statistically stable results. One out of the eight $70 \times 70$ samples we solved is shown in Fig. 4. To find this optimal configuration the program needed a little less than $16 \mathrm{hr}^{6}{ }^{6}$

At the beginning we considered $L$ up to 40 and $p=0.1,0.3,0.5,0.7$ and 0.9 . Note that, when $L$ is not too small, the value of the ground-state energy per spin for $p=0.1$ is very close to the value for $p=0.9$; the same holds also for the cases $p=0.3$ and $p=0.7$ (see Fig. 1). This is not surprising, since a (random) spin glass with concentration $p$ of negative interactions becomes (a random) one with a concentration $1-p$ when all couplings are multiplied by -1 . The following holds (and is proven in the appendix):

Proposition 1. The ground-state energies of a spin glass system with a concentration $p$ of negative interactions and the corresponding one with all couplings multiplied by -1 [concentration $(1-p)$ of negative interactions] are exactly the same if $L$ is even and differ by at most $4 L$ if $L$ is odd.

In other words, the deviations in the energy per spin are $O(1 / L)$, i.e., negligible for bigger sizes; however, these differences are noticeable for smaller samples.

Hence, for $L$ greater than 40 we just considered values of $p$ equal to $0.3,0.5$, and 0.9 . For each $L$ between 5 and 50 , we did $\left\lceil 20000 / L^{2}\right\rceil$ runs for every value of $p$. Figure 1 shows the averages for the ground-state energy for the different values of $p$ except for $p=0.5$. When $L$ is small the groundstate energy values for $p=0.1$ and $p=0.9$ show a very different behavior depending on whether $L$ is odd or even; these differences are not as high for the other values of $p$. For $p=0.1$ (or $p=0.9$ ) and $p=0.3$ ( or $p=0.7$ ) we estimate the ground-state energies for infinite systems as $E_{0.1,0.9}^{\infty} \approx-1.62$ and $E_{0.3,0.7}^{\infty} \approx-1.41$, respectively.

[^1]

Fig. 1. Ground-state energy for different concentrations.


Fig. 2. Ground-state energy for $p=0.5$.

In Fig. 2 we show the ground-state energies with $p=0.5$, which is the "classical" $\pm J$-spin glass. To make in particular the values for bigger lattices more accurate, we computed another 2000 ground states for the $p=0.5$ case. The error bars show the probable error $\sigma / \sqrt{N_{L}}$, where $N_{L}$ is the number of computed ground states. With a fitting function of the form $E_{v}(L)=E_{v}^{\infty}+c L^{-2}$, we find $E_{v}^{\infty}=-1.4022 \pm 0.0003$; with a function of the form $E_{c}(L)=E_{c}^{\infty}+c \exp (-a L)$, we get $E_{c}^{\infty}=-1.4007 \pm 0.0003$. Both results agree with the result given in ref. 11. Note that $L^{2}$ is the volume of the system; see ref. 14 for a brief discussion about both fitting functions.

In Fig. 3 the average CPU-times are shown. For better orientation there is also the graph of a function proportional to $L^{6}$ drawn as a straight line. Up to $L=50$ the CPU-time can be approximated by this polynomial; for bigger systems the variance is very high. It is not surprising that the running time can be bounded by a polynomial (in this size range), since branching was nearly never necessary $(99.53 \%$ of the samples were solved in the root node), so the program behaves mostly like a cutting plane algorithm with polynomial time separation routines.

Barahona and Titan ${ }^{(5)}$ reported computational results on $\pm J$ spinglass systems and published a certain $50 \times 50$ instance which took them


Fig. 3. CPU-time versus system size.


Fig. 4. Ground state of a sample with $L=70$. Positive couplings are drawn as solid lines.
about 17 hr CPU-time on a MIPS M-120 workstation using CPLEX as an LP-solver. We solved this one in 11 min .

Gropengiesser ${ }^{(11)}$ used biologically motivated algorithms for spin glasses with $p=0.5$. He approximated ground states up to grid sizes of $24 \times 24$ which took about 350 sec on a iPSC860 processor. Our program needed 17 sec on average to find an exact ground state for this sample size.

## 3. CONCLUSIONS

In this paper, we investigated exact ground-state computations for $\pm J$ Ising spin glasses on two-dimensional square lattices with nearest neighbor interactions and periodic boundary conditions. For the ground-state energy
of an infinite spin-glass system with $p=0.5$ we estimate $E_{0.5}^{\infty}=-1.4015 \pm$ 0.0008 . Since our approach is applicable also to three-dimensional systems, we plan to investigate them in the future.

## APPENDIX. PROOF OF PROPOSITION 1

In the proof we will use a result proven in ref. 9. For an edge-weighted graph $G=(V, E, w)$ with node set $V$, edge set $E$, and weight $w_{e}$ for each $e \in E$ and a subset $D \subseteq E$, let $w(D):=\sum_{e \in D} w_{c}$ and for a subset $S \subseteq V$ let $\delta(S):=\{(u, v) \in E \mid u \in S, v \notin S\}$ denote the cut induced by $S$.

Theorem 1. ${ }^{(9)}$ Let $G=(V, E, w)$ be an arbitrary weighted graph; let $\delta(S)$ denote an arbitrary cut of $G$. The weighted graph $G_{\delta(S)}=\left(V, E, w^{\prime}\right)$ is the graph that is obtained by switching the weights of the graph $G$ along the cut $\delta(S)$ i.e.,

$$
w_{c}^{\prime}=\left\{\begin{array}{lll}
-w_{e} & \text { if } \quad e \in \delta(S) \\
w_{c} & \text { if } \quad e \notin \delta(S)
\end{array}\right.
$$

If $W_{G}$ is the weight of a maximum cut in $G$, then the weight $W_{G_{\bar{\alpha}, 5}}$ of a maximum cut in $G_{\delta(S)}$ is given by

$$
W_{G_{u(S)}}=W_{G}-w(\delta(S))
$$

Theorem 2. Let $G=(V, E, w)$ be a toroidal grid graph with $L \times L$ nodes and weights $w_{c}= \pm 1$ on the edges. Let $G^{\prime}=\left(V, E, w^{\prime}\right)$ be the graph that is obtained by multiplying the weight of every edge in $G$ by -1 . The weight $W_{G}$, of a maximum cut in $G^{\prime}$ is given by

$$
W_{G^{\prime}}=\left\{\begin{array}{ll}
W_{G}-w(E) & \text { if } L \text { is even } \\
W_{G}-w(E)+\Delta & \text { with }|\Delta| \leqslant 2 L
\end{array} \quad \text { if } L\right. \text { is odd }
$$

Proof. $L$ is even. Since $E$ is a cut in $G$, the graph $G^{\prime}$ can be obtained by switching along $E$. Applying Theorem 1 directly results in

$$
W_{G^{\prime}}=W_{G}-w(E)
$$

$L$ is odd. Consider the graph $\tilde{G}=(\tilde{V}, \tilde{E}, \tilde{w})$ which is obtained from $G$ by adding one row and one column of nodes and edges (see Fig. 5). All new edges have weight 0 . Because $\tilde{G}$ is again a toroidal grid graph with an even $\tilde{L}$ now, Theorem 1 can be applied to the graphs $\tilde{G}$ and $\tilde{G}^{\prime}$ obtained from $\tilde{G}$ by multiplying all edge weights by -1 . So we have

$$
\begin{equation*}
W_{\bar{G}^{\prime}}=W_{\tilde{G}^{-}}-w(E) \tag{Al}
\end{equation*}
$$

because $\tilde{w}(\tilde{E})=w(E)$.


Fig. 5. Converting $G$ into $\tilde{G}$ and $S$ into $\tilde{S}$.
We can easily extend any $S \subseteq V$ to $\tilde{S}$, where $S \subseteq \tilde{S} \subseteq \tilde{V}$ in order to obtain a cut in $\tilde{G}$ of the same weight $\tilde{w}(\delta(\tilde{S}))=w(\delta(S))$ as the cut in $G$ induced by $S$. Namely, let $\tilde{s} \in \tilde{V} \backslash V$ be adjacent to the nodes $u$ and $v$, where $u, v \in V$. (All but one node in $\tilde{V} \backslash V$ are of that kind.) Let $\tilde{e}=(v, \tilde{s}) \in \tilde{E} \backslash E$, $\tilde{e}^{\prime}=(u, \tilde{s}) \in \tilde{E} \backslash E$, and $e=(u, v) \in E$, such that $\tilde{w}(\tilde{e})=w(e)$ and $\tilde{w}\left(\tilde{e}^{\prime}\right)=0$. Then let $\tilde{s} \in \tilde{S}$ if and only if $u \in S$, because then the edge $\tilde{e}$ is in the cut $\delta(\tilde{S})$ if and only if $e \in \delta(S)$. The same reasoning applies to $G^{\prime}$ and $\widetilde{G}^{\prime}$, so we have

$$
\begin{equation*}
W_{G} \leqslant W_{\widetilde{G}} \quad \text { and } \quad W_{G^{\prime}} \leqslant W_{\widetilde{G}^{\prime}} \tag{A2}
\end{equation*}
$$

Let $\delta(\tilde{S})$ with $\tilde{S} \subseteq \tilde{V}$ be a cut of maximum weight in $\widetilde{G}$, i.e., $W_{\tilde{G}}=\tilde{w}(\delta(\tilde{S}))$. Then we can restrict $\tilde{S}$ to $S=\tilde{S} \cap V$ in order to obtain a cut $\delta(S)$ in $G$ whose weight is at least $W_{\tilde{G}}-2 L$ because there are $2 L$ edges in $E$ that are connected to nodes in $\tilde{V} \backslash V$ and the edges connecting only nodes in $\tilde{V} \backslash V$ have all weight zero by construction. Since the value of this cut in $G$ is in turn a lower bound for the value of an optimal cut in $G$, we obtain

$$
\begin{equation*}
W_{G} \geqslant W_{\bar{G}}-2 L \quad \text { and } \quad W_{G^{\prime}} \geqslant W_{\tilde{G}^{\prime}}-2 L \tag{A3}
\end{equation*}
$$

because the same holds for $G^{\prime}$ and $\tilde{G}^{\prime}$.
Combining (A1)-(A3), we have

$$
\begin{aligned}
& W_{G}-w(E)-2 L \stackrel{\left(A_{2}\right)}{\leqslant} W_{\tilde{G}}-w(E)-2 L \stackrel{\left(A_{1}\right)}{=} W_{G^{\prime}}-2 L \\
& \stackrel{\left(A^{3}\right)}{\leqslant} W_{G^{\prime}} \\
& \stackrel{(\mathrm{A} 2)}{\leqslant} W_{\tilde{G}^{\prime}} \stackrel{(\mathrm{A} 1)}{=} W_{\tilde{G}}-w(E) \\
& \stackrel{\left(\mathrm{A}^{3}\right)}{\leqslant} W_{G}-w(E)+2 L \quad \mathrm{QED}
\end{aligned}
$$

Since the energies of configurations in spin-glass systems represented by graphs $G=(V, E, w)$ and $G^{\prime}=\left(V, E, w^{\prime}\right)$ are given by

$$
E=-2 W_{G}+w(E) \quad \text { and } \quad E^{\prime}=-2 W_{G^{\prime}}+w^{\prime}(E)
$$

respectively, ${ }^{(8)}$ and

$$
w^{\prime}(E)=-w(E)
$$

the ground-state energies of a spin-glass system with concentration $p$ of negative interaction and the corresponding one with all couplings multiplied by -1 [concentration $(1-p)$ ] are exactly the same if $L$ is even and differ by at most $4 L$ if $L$ is odd. So the energy per spin difference is at most $4 L / L^{2}=4 / L$ for odd $L$.

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[^0]:    ${ }^{\prime}$ Institut für Informatik, Universität zu Köln, 50969 Köln, Germany.
    ${ }^{2}$ On leave from IASI-CNR, 00185 Rome, Italy.
    ${ }^{3}$ MPI für Informatik, 66123 Saarbrücken, Germany.
    ${ }^{4}$ Institut für Angewandte Mathematik, Universität Heidelberg, 69120 Heidelberg, Germany.
    ${ }^{5}$ IASI-CNR, 00185 Rome, Italy.

[^1]:    ${ }^{6}$ An electronic version of each sample is available from the authors; e-mail: diehl@infor-matik.uni-koeln.de.

